CSCE 463/612
Networks and Distributed Processing
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Application Layer V
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Homework #2

- Unlike HTTP, all fields are binary
  - Make sure to refresh pointer usage
- Question format:
  - Fill in the values (flags: DNS_QUERY and DNS_RD, nQuestions = 1)
  - Allocate memory for the packet
  - Write question into buffer

```
class QueryHeader {
  u_short type;
  u_short class;
};
```

```
class FixedDNSHeader {
  u_short ID;
  u_short flags;
  u_short questions;
...;
};
```
Homework #2

High-level operation for DNS questions:

```c
char packet [MAX_DNS_LEN]; // 512 bytes is max
char host[] = "www.google.com";
int pkt_size = strlen(host) + 2 + sizeof(FixedDNSHeader) + sizeof(QueryHeader);

// fixed field initialization
FixedDNSHeader *dh = (FixedDNSHeader *) packet;
QueryHeader *qh = (QueryHeader*) (packet + pkt_size - sizeof(QueryHeader));
dh->ID = ...
dh->flags = ...
...
qh->type = ...
qh->class = ...

// fill in the question
MakeDNSquestion (dh + 1, host);
// transmit to Winsock
sendto (sock, packet, ...);
```

If packet is incorrectly formatted, you will usually get no response; use Wireshark to check outgoing packets
Homework #2

• Formation of questions:
  makeDNSquestion (char* buf, char *host) {
    while(words left to copy){
      buf[i++] = size_of_next_word;
      memcpy (buf+i, next_word, size_of_next_word);
      i += size_of_next_word;
    }
    buf[i] = 0; // last word NULL-terminated
  }

• All answers start with an RR name, followed by a fixed DNS reply header, followed by the answer
  - Uncompressed answer (not common)
    0x3 “irl” 0x2 “cs” 0x4 “tamu” 0x3 “edu” 0x00 <DNSanswerHdr> <ANSWER>
  - Compressed (2 upper bits 11, next 14 bits jump offset)
    0xC0 0x0C <DNSanswerHdr> <ANSWER>

• For type-A questions, the answer is a 4-byte IP
**Homework #2**

- To check the header
  - Hex printout on screen
  - Use Wireshark
- **What is `sizeof(DNSanswerHdr)`?**
  - The actual size is 10 bytes, but most structures in C/C++ are aligned/padded to 4 (or 8) byte boundary
- Remember to change struct packing of all classes that define binary headers to 1 byte
- Caveats (must be properly handled):
  - Exceeding array boundaries on jumps
  - Infinite looping on compressed answers

```c
class DNSanswerHdr {
    u_short type;
    u_short class;
    u_int ttl;
    u_short len;
};
```

```c
#pragma pack(push,1) // define class here
#pragma pack(pop)
```
How to check if compressed and read 14-bit offset?
- Suppose array `char *ans` contains the reply packet
- The answer begins within this array at position `curPos`

```
1 1 x x x x x x x x x x x x x
```

14 bits

```
char *ans; // points to reply buffer
if (ans[curPos] >= 0xC0)
   // compressed; so jump
else
   // uncompressed, read next word
```

```
char *ans; // points to reply buffer
if ( (ans[curPos] >> 6) == 3)
   // compressed; so jump
else
   // uncompressed, read next word
```

// computing the jump offset
int off = ((ans[curPos] & 0x3F) << 8) + ans[curPos + 1];

The top two algorithms will generally fail
- Use only `unsigned` chars when reading buffer!
Note that jumps may appear in mid-answer 0x3 “irl” 0xC0 0x22 <DNSanswerHdr> <ANSWER>

Jumps may be nested, but must eventually end with a 0-length word
  - Need to remember the position following the very first jump so that you can come back to read the answer

Replies may be malicious or malformatted
  - Homework must avoid crashing

If AAAA (IPv6) answers are present, skip
  - Use DNSanswerHdr::len to skip unknown types

Caution with SuddenLink and dorms
  - Too many malformed packets may raise an alarm
Chapter 2: Roadmap

2.1 Principles of network applications
2.2 Web and HTTP
2.3 FTP
2.4 Electronic Mail
   - SMTP, POP3, IMAP
2.5 DNS (extras)
2.6 P2P file sharing
Domain Flux

- Viruses, trojan horses, rootkits, and various malware affect millions of computers today
- 20 years ago, viruses mostly performed pranks or corrupted data, but this has changed
  - Modern attacks are often driven by financial gains
- Infected hosts are organized into botnets
  - Large collection of computers under control of a botmaster
- Early botnets used IRC (Internet Relay Chat) to send and receive commands

Case at hand: Jeremy Jaynes (arrested in 2004) sent an estimated 10M emails/day, pulling $750K/month income
• Eventually, ISPs started blocking IRC traffic
  – Also, IRC servers were easy targets for shutdown and filtering (e.g., detection of encrypted commands and botnet channels)
• New generation of botnets uses dynamically changing rendezvous points called **C&C** (command & control)
  – Stealthy because C&C’s IP can rapidly change over time
  – Main problem: how does the botnet find the current C&C?
**Domain Flux 3**

- **Fast flux** is a method for discovering the IP address of C&C and other resources the botnet may need
  - Botmaster registers a domain (say xyz.com) and controls the DNS server ns.xyz.com
- Botnet contacts nameserver ns.xyz.com and obtains the current IP of the C&C (or multiple ones)
  - Performs a type-A lookup on hash.xyz.com
### Domain Flux 4

- Main defense against botnet traffic is blocking communication with the C&C
  - Fast Flux makes it harder since the C&C changes over time and is load-balanced across several hosts
  - When C&C is blocked, botnet learns other locations quickly

- Fast flux can also be used to serve phishing content
  - Suppose email arrives to user with a link to www77.xyz.com
  - Botnet uses DNS to serve this request from a variety of compromised hosts
Several benefits to serving HTTP content using fast flux
- Difficult to trace IPs hosting content or block malicious URLs
- Botnet is failure resilient -- if hosts are cleaned or go offline, there is automatic fail-over to other live hosts
- Cheap in terms of bandwidth, simple to implement

However, there is a problem
- Suppose ISP or SpamAssasin blocks all requests for xyz.com or registrar disables xyz.com?
- If xyz.com is taken down, the botnet freezes

Domain flux aims to solve this issue
- Botnet constantly generates random domain names and tries to resolve them to find the C&C
- Much more difficult to trace and block

Nowadays, TLD servers auto-detect fastflux and block suspected domains in conjunction with the registrar
Domain Flux 6

• Toy example:
  – Suppose botnet computers generate names using this sequence: 1.com, 2.com, 3.com, 5.com, 8.com, 13.com, etc.
  – Current domain name stays in effect until it is blocked
  – Initially, botmaster registers 1.com and 34.com
  – When 1.com gets blocked, the botnet automatically switches to 34.com, while botmaster registers 144.com, and so on

• In reality, the botnet goes through thousands of failed lookup attempts until it finds an active domain
  – Can be detected from a huge number of failed DNS queries

• Domains may be too random to be human-produced
  – If so, machine-learning algorithms can be used to detect infected hosts that are attempting domain flux
Domain Flux 7

• In some cases, reverse engineering the random generator allows one to predict future domain names
  - By registering these domains, botnets can be hijacked
  - Researchers have shown this is possible in B. Stone-Gross et al., “Your botnet is my botnet: Analysis of a botnet takeover,” ACM CCS, 2009.

• How large are botnets? Some examples:
  - BredoLab (2009): 30M hosts, 3.6B emails/day
  - Conficker (2008): 10.5M hosts, 10B emails/day
  - Cutwail (2007): 1.5M hosts, 74B emails/day
  - Torpig (paper above): 180K hosts (theft of 500K bank accounts, credit cards)
  - Avalanche (2008-2016): phishing botnet w/500K hosts
Chapter 2: Roadmap

2.1 Principles of network applications
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2.5 DNS
2.6 P2P file sharing
Hybrid P2P

- Napster (1999)
  - Application-layer protocol over TCP
  - Centralized directory server

- Sequence of steps
  - Connect to server, login
  - Upload your IP/port + list of files
  - Give server keywords for search
  - Select “best” answer (ping)
  - Download from peer

- Single point of failure
- Performance bottleneck
- Target for litigation due to copyright infringement
Decentralized P2P

- Napster folded in 2002
  - Other P2P systems took over (Gnutella, KaZaA, BitTorrent, eDonkey)
- Gnutella/0.4 (2001)
  - Public-domain protocol
  - Fully distributed design
- Many Gnutella clients implementing protocol
  - Limewire, Morpheus, BearShare

- How to find content?
  - Idea: construct a graph
    - Edge between peer X and Y iff there’s a TCP connection between them
- All active peers and edges are called an overlay network
  - Peer typically connected to < 30 neighbors
- Search proceeds by flooding up to some depth
  - Limited-scope flooding
Decentralized P2P

- Queries are P2P
  - Inefficient due to huge volumes of traffic
  - Average degree k, depth of flood d, overhead \((k-1)^d\)

- Downloads are P2P from a single user
  - Unreliable (peer departure or failure kills transfer)
  - Inefficient (asymmetry of upstream/downstream bandwidth)

- Join protocol (bootstrapping)
  - Find an entry peer X, flood its neighbors to obtain more candidates, establish connections to those who accept
Hierarchical P2P

- Gnutella/0.4 scaled to about 25K users and then choked
- Alternative construction proposed by KaZaA (2002)
  - Peer is either a group leader (supernode) or assigned to one
- Group leader tracks the content of all its children, acting like a mini-Napster
  - Peers query their group leaders, which flood the supernode graph until some number of matches found
  - Query-hits not routed, but sent directly to original supernode
Hierarchical P2P

- With 150 neighbors, this architecture is 150 times more efficient than Gnutella/0.4 in message overhead
  - With 389M downloads as of 2008, KaZaA was more popular than Napster ever was, accounting for 50% of ISP bandwidth in some regions and running 3M concurrent users
- Gnutella/0.6 soon adopted the same structure
  - Scaled to 6.5M online users, 60M unique visitors per week
- Additional features
  - Hashed file contents to identify exact version of files
  - Upload and request queuing at each user, rate-limiting
  - Parallel downloads from multiple peers
  - Support for crawl requests that reveal neighbors
Other P2P

• **Terminology**: user holding a complete file is a **seed**
  - Traditional systems download only from seeds
  - Seed departs, transfer fails

• **Idea**: let non-seeds grab chunks from each other
  - Peers organize into a group (torrent) based on the file they’re downloading

• Traditional systems download files **sequentially**
  - Starvation for final blocks

• **Idea**: maximize availability
  - Participants forced to serve chunks they have to others
  - Rarest chunk in torrent is always replicated first

• **Known as BitTorrent (2001)**
  - Protocol with many implementations
  - Requires **trackers** to keep torrent membership
  - At any time, more concurrent users that YouTube and Facebook combined

• **Built-in incentives to share**
  - Rate-limiting (choking) based on upload activity
**Other P2P**

- **Tor (Onion Router)**
  - Anonymity network of peers
  - Each packet sent through a random chain of P2P nodes
    - Final user relays packet towards destination
    - Return packets processed similarly along reverse path
- **Tor can be run thru an API**
  - Extremely slow
  - Many exit points are known and blocked by Google
- **Roughly 36M users**
- **Freenet**
  - Anonymous information exchange, hiding identities of communicating parties
- **Skype chat**
  - Video streaming services either directly between users or relayed through non-firewalled peers
- **Distributed Hash Tables**
  - General class of P2P systems that map information into high-dimensional search space with guaranteed log(N) bounds on delay to find content