CSCE 313-200
Introduction to Computer Systems
Spring 2022

Practice III
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Quiz 3

- Write proper synchronization for a train tunnel

```c
Train
TryEnteringTunnel (int dir) {
  mutex[dir].Lock();
  if (trains[dir]++ == 0)
    occupied.Wait();
  mutex[dir].Unlock();

  semaMaxN.Wait();
  PassThruTunnel(x, dir);
  semaMaxN.Release();

  mutex[dir].Lock();
  if (--trains[dir] == 0)
    occupied.Release();
  mutex[dir].Unlock();
}
```
• Print spooler system
  - Main rule: combined size of Q1 and Q2 cannot exceed M
• Version #1: without the combined max, each queue has an independent size limit

**Quiz3**

```
P1
x = ObtainItem();
semaEmptyQ1.Wait();
m1.Lock();
Q1.push(x);
m1.Unlock();
semaFullQ1.Release();
```

```
P2
semaFullQ1.Wait();
m1.Lock();
y = Q1.pop();
m1.Unlock();
semaEmptyQ1.Release();
z = Process (y);
semaEmptyQ2.Wait();
m2.Lock();
Q2.push(z);
m2.Unlock();
semaFullQ2.Release();
```

```
P3
semaFullQ2.Wait();
m2.Lock();
w = Q2.pop();
m2.Unlock();
semaEmptyQ2.Release();
ProcessAndDiscard (w);
```
• Version #2: with the max, but deadlock-prone

Semaphore disk = {M,M};

P1

x = ObtainItem();
disk.Wait();
m1.Lock();
Q1.push(x);
m1.Unlock();
semaFullQ1.Release();

P2

semaFullQ1.Wait();
m1.Lock();
y = Q1.pop();
m1.Unlock();
disk.Release();
z = Process (y);  
disk.Wait();
m2.Lock();
Q2.push(z);
m2.Unlock();
semaFullQ2.Release();

P3

semaFullQ2.Wait();
m2.Lock();
w = Q2.pop();
m2.Unlock();
disk.Release();
ProcessAndDiscard (w);
Quiz3

• Version #3: do not release disk semaphore in P2

Semaphore disk = \{M,M\};

\begin{align*}
\text{P1} \\
x &= \text{ObtainItem();} \\
disk.\text{Wait();} \\
m1.\text{Lock();} \\
Q1.\text{push(x);} \\
m1.\text{Unlock();} \\
\text{semaFullQ1.Release();}
\end{align*}

\begin{align*}
\text{P2} \\
\text{semaFullQ1.Wait();} \\
m1.\text{Lock();} \\
y &= Q1.\text{pop();} \\
m1.\text{Unlock();} \\
// \text{remove disk.Release();} \\
z &= \text{Process (y);} \\
// \text{remove disk.Wait();} \\
\text{semaFullQ1.Release();}
\end{align*}

\begin{align*}
\text{P3} \\
\text{semaFullQ2.Wait();} \\
m2.\text{Lock();} \\
w &= Q2.\text{pop();} \\
m2.\text{Unlock();} \\
disk.\text{Release();} \\
\text{ProcessAndDiscard (w);} \\
\end{align*}

• What if P2 makes K items for each extracted from Q1?
• Assume N processes sharing M resources
  - Process i eventually wants to hold $W_i$ resources
  - Resources are obtained non-atomically
  - After getting all of its resources, process releases them
• Maximum # of resources $R$ that still lead to deadlock?
  - Suppose $W_1 = 6$, $W_2 = 3$, $W_3 = 14$
  - Then $M > R$ guarantees no deadlock and $M = R$ allows one
• Writing:
  \[ R = \sum_{i=1}^{N} (W_i - 1) < M \]
  - we obtain:
  \[ \sum_{i=1}^{N} W_i - N < M \quad \Rightarrow \quad \sum_{i=1}^{N} W_i < M + N \]
String Search

• How fast is homework #3 with 200K keywords?
  - Roughly 9.1 KB/s, 38 days to parse the big file
• Using all 8M unique words in large Wikipedia?
  - Speed 240 bytes/s, roughly 4 years to finish (using 12 cores)
• Focus of computer science has always been efficiency
  - Quicksort vs bubble sort, hashing vs sorting, binary vs linear search, min-heap vs linear min()
  - Substring search is another example
• Start with single-string search
  - Assume some text and a given keyword
  - Need to find all occurrences of keyword in text
  - Matches do not have to be complete words
• Naïve method #1: use strcmp or memcmp
• Naïve method #2: use strstr
  - Runs somewhat faster, but still far from optimal
• Example of method #1:
  - Worst-case complexity?
  - N = length of text, M = word size, then (N-M)*M

while (off < bufSize - wordLen) {
    if (memcmp (buf + off, word, wordLen) == 0)
        found ++;
    off ++;
}

char *match = buf;
buf [bufSize] = 0;
while (true) {
    match = strstr (match, word);
    if (match == NULL)
        break;
    found ++;
    match ++;
}
Single String

<table>
<thead>
<tr>
<th>text</th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>Q</th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>D</th>
<th>A</th>
<th>B</th>
<th>Z</th>
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<tbody>
<tr>
<td>miss</td>
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<tr>
<td>word</td>
<td>A</td>
<td>B</td>
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Step 2

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<tr>
<th>text</th>
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<th>B</th>
<th>C</th>
<th>Q</th>
<th>A</th>
<th>B</th>
<th>C</th>
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Step 3

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</table>

Step 4

<table>
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<th>B</th>
<th>C</th>
<th>Q</th>
<th>A</th>
<th>B</th>
<th>C</th>
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</table>

...
**Single String**

- Naïve takes 7 comparisons to move 4 bytes
  - Total complexity of getting past 12 bytes is 23 comparisons
- Knuth-Morris-Pratt (KMP), 1977:

  **Step 1**

<table>
<thead>
<tr>
<th>text</th>
<th>A</th>
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<th>C</th>
<th>Q</th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>D</th>
<th>A</th>
<th>B</th>
<th>Z</th>
<th>D</th>
<th>...</th>
</tr>
</thead>
<tbody>
<tr>
<td>word</td>
<td>A</td>
<td>B</td>
<td>C</td>
<td>D</td>
<td>A</td>
<td>B</td>
<td>D</td>
<td>Q</td>
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</table>
  
  **Step 2**

<table>
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<td>Q</td>
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</tr>
</tbody>
</table>
Single String

text: A B C Q A B C D A B Z D ...
word: A B C D A B D Q ...
step 3

A B C Q A B C D A B Z D ...
miss

word: already matched

A B C Q A B C D A B Z D ...
miss

step 4

A B C D A B D Q ...

A B C Q A B C D A B Z D ...
miss

step 5

A B C D
Single String

- Total 6 steps, 15 comparisons to pass 12 bytes
- How does it work?
  - Each character needs two lookup tables (LUTs) – by how many bytes to move after a non-match in this position and where in the word to re-start on the next attempt

<table>
<thead>
<tr>
<th>word</th>
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<th>D</th>
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<th>B</th>
<th>D</th>
<th>Q</th>
</tr>
</thead>
<tbody>
<tr>
<td>move</td>
<td>1</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>5</td>
<td>5</td>
<td>4</td>
<td>7</td>
</tr>
<tr>
<td>re-start</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>0</td>
</tr>
</tbody>
</table>

tables built offline, fit in L1 cache
**Single String**

- **Boyer-Moore** (BM), 1977:
  - Uses not just distance, but also the mismatched character
- Matching goes right to left, until a mismatch
  - Off is examined position in text

<table>
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<th>B</th>
<th>C</th>
<th>C</th>
<th>Q</th>
<th>B</th>
<th>Z</th>
<th>D</th>
<th>Q</th>
<th>1</th>
<th>2</th>
</tr>
</thead>
<tbody>
<tr>
<td>off</td>
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</tbody>
</table>

- After a miss, two hash tables move the word forward:
  - Slide[dist]: based on the # of matched characters
  - Shift[char]: based on mismatched character text[off]
Single String

- In the example above
  - Mismatch distance is 0, so slide by 1 char
  - Mismatch char = C, so shift by 5
- After moving off by 5:

  ![Diagram showing text and word comparison]

  - In this case, mismatch occurs at text[off] = Z:
    - Mismatch distance = 2, slide word by 8
    - Mismatch char = Z, shift word by 6
For words that have rare letter combinations, we can be skipping by M each time
  - Best case complexity is sub-linear, i.e., N/M comparisons

Typically faster than KMP for larger M
Can we do better?
Notice that BM gets stuck on popular characters, while ideally it should skip most examined locations
  - E.g., “zebra” incurs detailed inspection any time it hits an ‘e’
Idea: set up a hash table with 2-byte combinations
  - E.g., “ze”, “eb”, “br”, “ra” which are much more rare
  - Then scan the text using an unsigned short (2-byte) pointer
Caveat: don’t know alignment of the word, may hit something like “_z” and miss the word
  - Need to set up wildcard entries ?z and a? for all possible leading and trailing characters
  - If only full words are needed, ? will be a white space
Why was homework #3 so inefficient?

Idea: do not compare current byte to all strings, only to those that can potentially be a match

Rabin-Karp (RK), 1987
- Assume $M$ is the smallest keyword length
- Compute a hash $H$ of the next $M$ chars from current location
- Hit a hash table, compare with words that tie for that hash
- Speed is only based on the length of collision chains
Multiple Strings

• After hash table lookup, slide by one byte forward, recompute the hash of the next M chars

• Notice that M-1 chars are the same in both hashes
  - Main twist of the algorithm is to use a rolling hash, which obtains $H_{i+1}$ from $H_i$ in $O(1)$ time

• Treating hashes as base-B integers, we have
  - $H_0 = \text{str}[0] \times B^{M-1} + \text{str}[1] \times B^{M-2} + \ldots + \text{str}[M-1]$
  - $H_{i+1} = (H_i \times B + \text{str}[i+M]) \mod B^M$
Wrap-up

- Larger M means fewer collisions and faster operation
- With M = 3 and 216K strings, RK runs at 20MB/s
  - 2000 times faster than the naïve method
- Indexing a file with unknown keywords is slightly different, but the idea is similar to RK
  - Homework #4 explores this in more detail
- Main goal is to design code that processes all 4.5B words in large Wikipedia in ~35 sec (135M wps)
  - 3.7M times faster than the method in homework #3
- Homework #4 has 3 checkpoints
  - The first two should be done early
  - Checkpoint #3 is more complex, uses virtual memory